



DRONACHARYA
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INTELLIGENT SYSTEMS (CSE-303-F)

Section A

Alpha Beta Pruning

Motivations

- Breadth-first, depth-first, hill-climbing, best-first, and A* assume a non-hostile search space.
- The goals just sit there somewhere in the graph.
- The goals do not try to elude you.
- The 8-puzzle game did not try to stop you from reaching the goal.
- Your tic-tac-toe opponents reacted to your moves randomly.
- But in a real 2-person game, your opponent does try to beat you and make it difficult for you to reach your goal.
- Minimax search can be applied in an adversarial search space.
- Alpha-beta pruning can be used to cut bad branches (moves) in the game tree to improve minimax search time.

Objectives

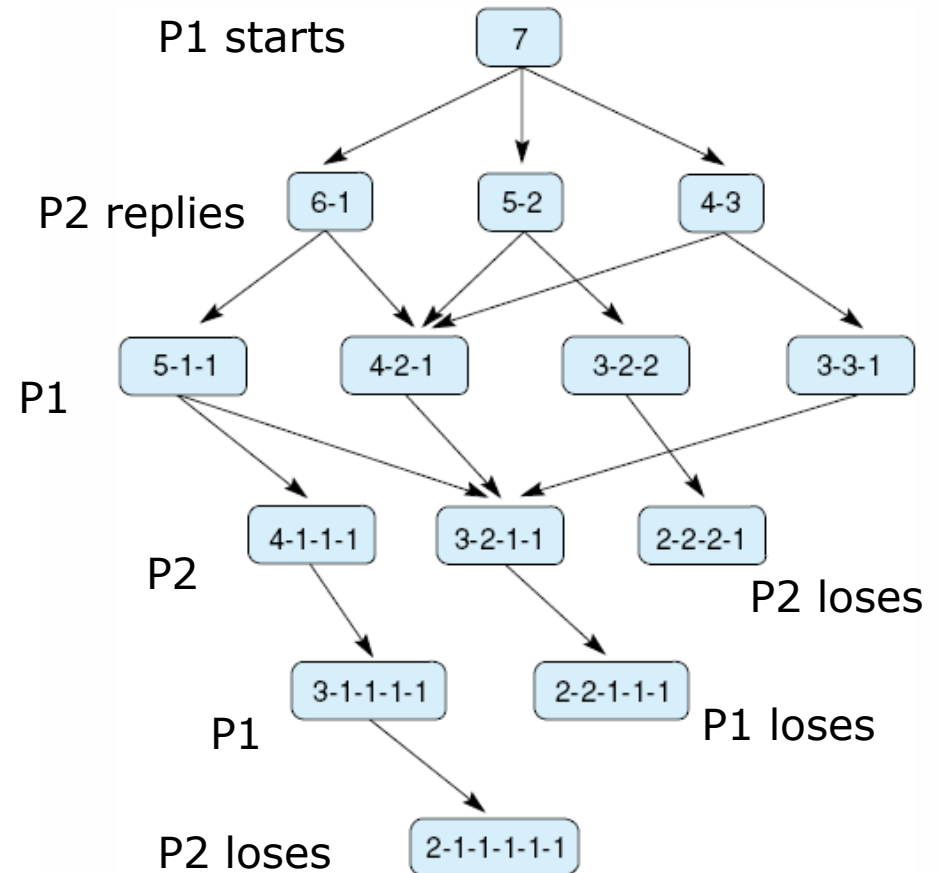
1. Adversarial search space: MAX vs. MIN
2. A simple game tree: Nim-7
3. Minimax on Nim-7
4. Minimax on tic-tac-toe looking 3 plies ahead
5. Alpha-beta pruning

Two people games

- Solved games
 - Tic-tac-toe
 - Four In A Line
 - Checkers
- Impressive games played by robots
 - Othello bot is much stronger than any human player
 - Computer chess beat the human world champions
 - TD-Gammon ranked among top 3 players in the backgammon world
- Future bot challenges to humans
 - Poker bots play respectfully at world-class level
 - Computer bridge programs play competitively at national level
 - Go bots are getting more serious in the amateur ranking

Complete game tree for Nim-7

- 7 coins are placed on a table between the two opponents
- A move consists of dividing a pile of coins into two nonempty piles of different sizes
- For example, 6 coins can be divided into piles of 5 and 1 or 4 and 2, but not 3 and 3
- The first player who can no longer make a move loses the game



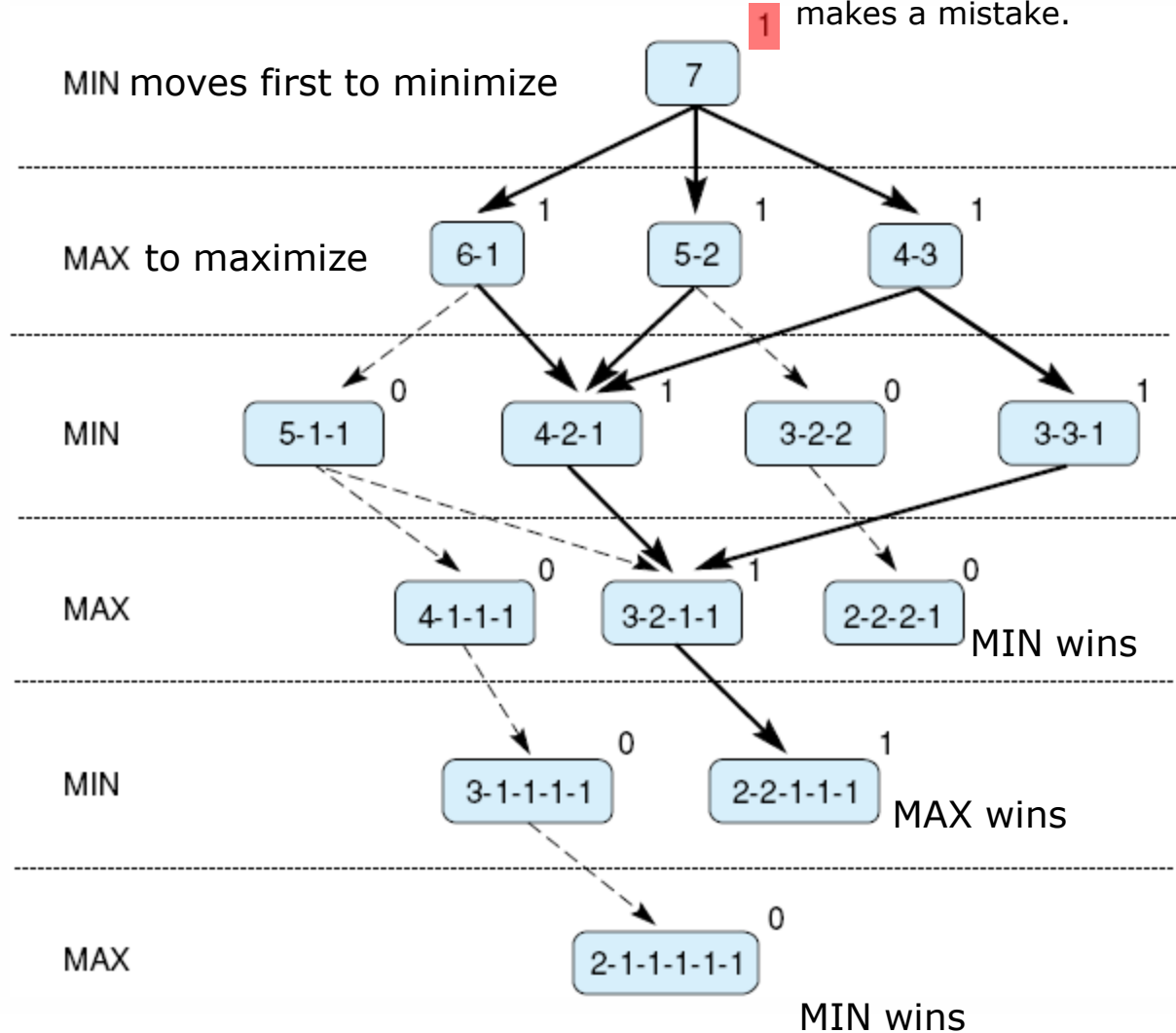
MIN vs. MAX in a Nim game

The best that MIN (Player1) can do is to lose unless Player2 makes a mistake.

Node score = 0 means MIN wins.

1 means MAX wins.

Bold edges indicate forced win for MAX, Player2.

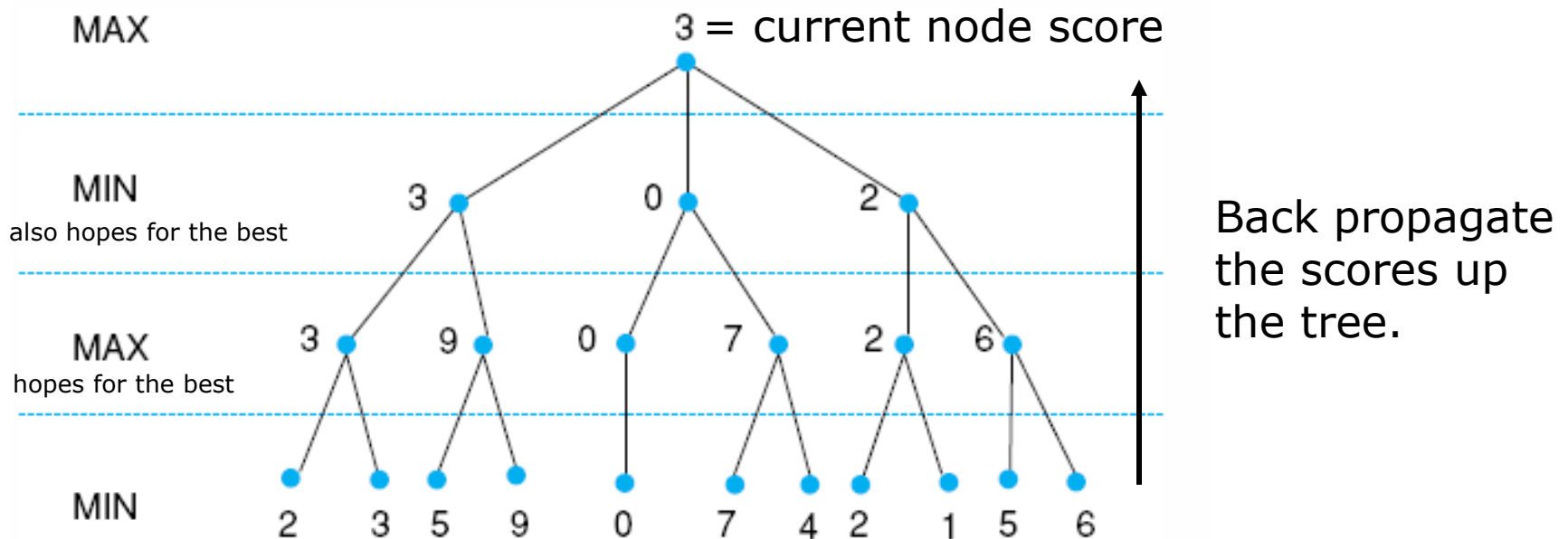


Minimax to fixed ply depth

- Instead of Nim-7, imagine the chess game tree.
- Chess game tree is too deep.
 - cannot expand the current node to terminating (leaf) nodes for checkmate.
- Use fixed ply depth look-ahead.
 - Search from current position to all possible positions that are, e.g., 3-ply ahead.
- Use heuristic to evaluate all these future positions.
 - $P=1, N=B=3, R=5, Q=9$
 - Assign certain weight to certain features of the position (dominance of the center, mobility of the queen, etc.)
 - summarize these factors into a single number.
- Then propagating the scores back to the current node.

MAX calculates the current node score

Look 3 plies ahead.



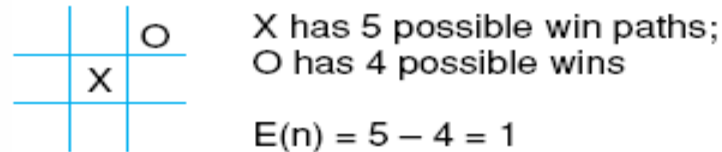
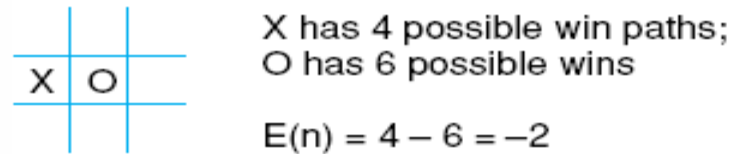
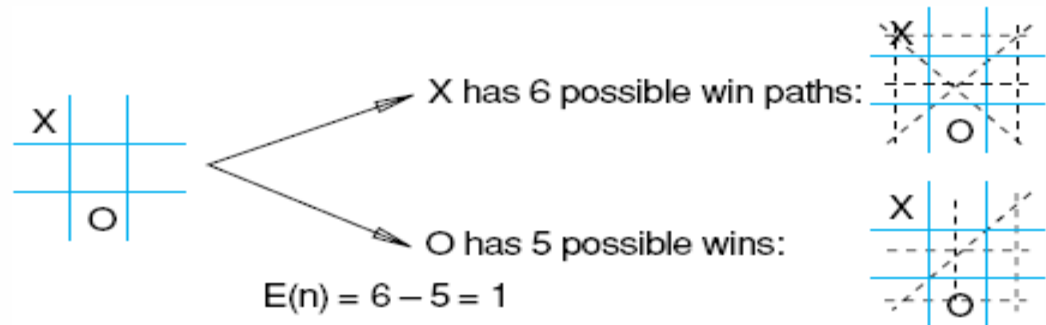
Use heuristic $h(n)$ for each of these future positions.

A stronger heuristic will beat a weaker heuristic.

A farther look-ahead will beat a near-sighted look-ahead.

Computer chess routinely uses complex heuristics analyzing material and positional advantages and looks 40 plies ahead.

Heuristic measuring for adversarial tic-tac-toe



Heuristic is $E(n) = M(n) - O(n)$

where $M(n)$ is the total of My possible winning lines

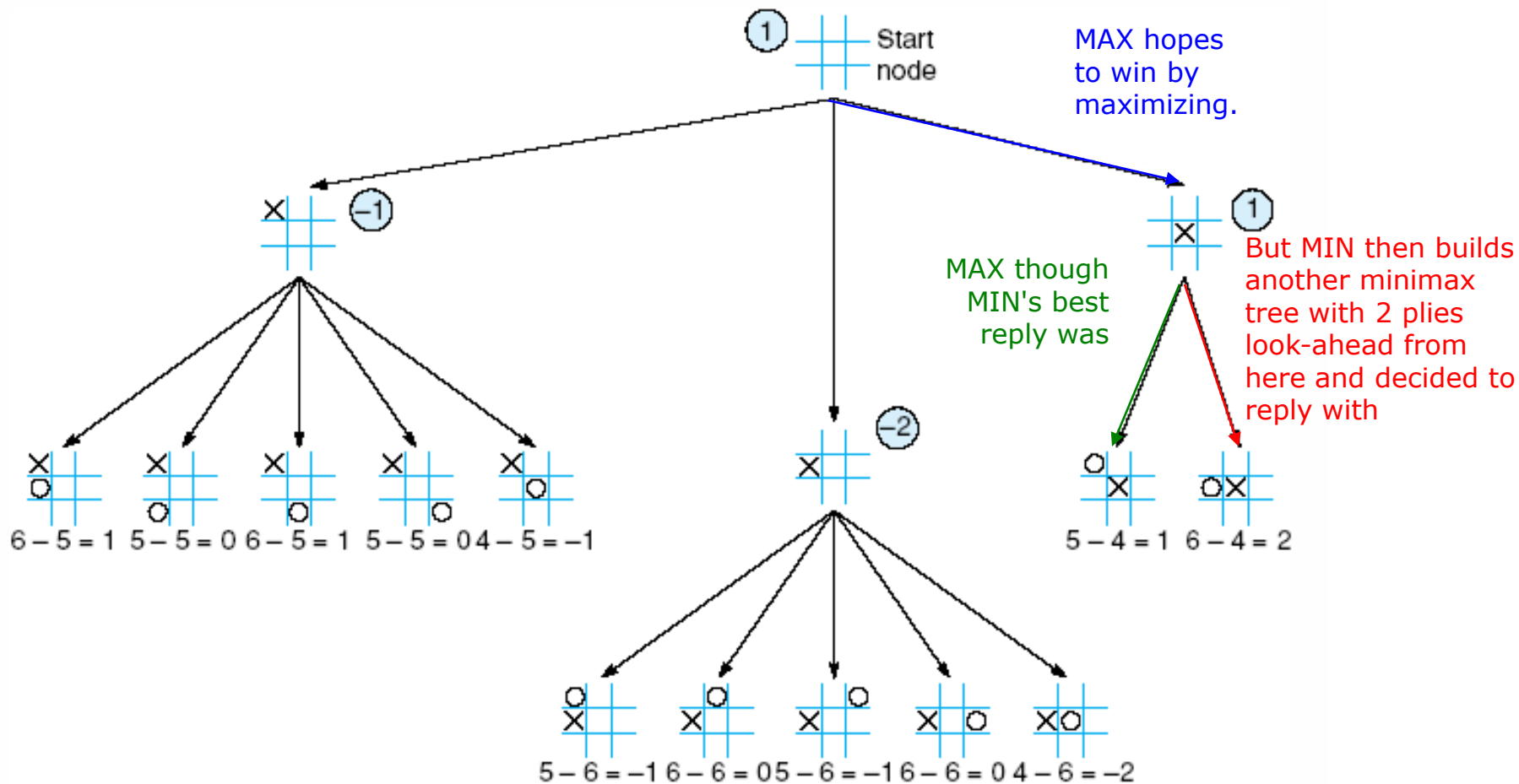
$O(n)$ is total of Opponent's possible winning lines

$E(n)$ is the total Evaluation for state n

Maximize $E(n)$

$E(n) = 0$ when my opponent and I have equal number of possibilities.

MAX makes his first move



MAX's 2nd move: look ahead analysis

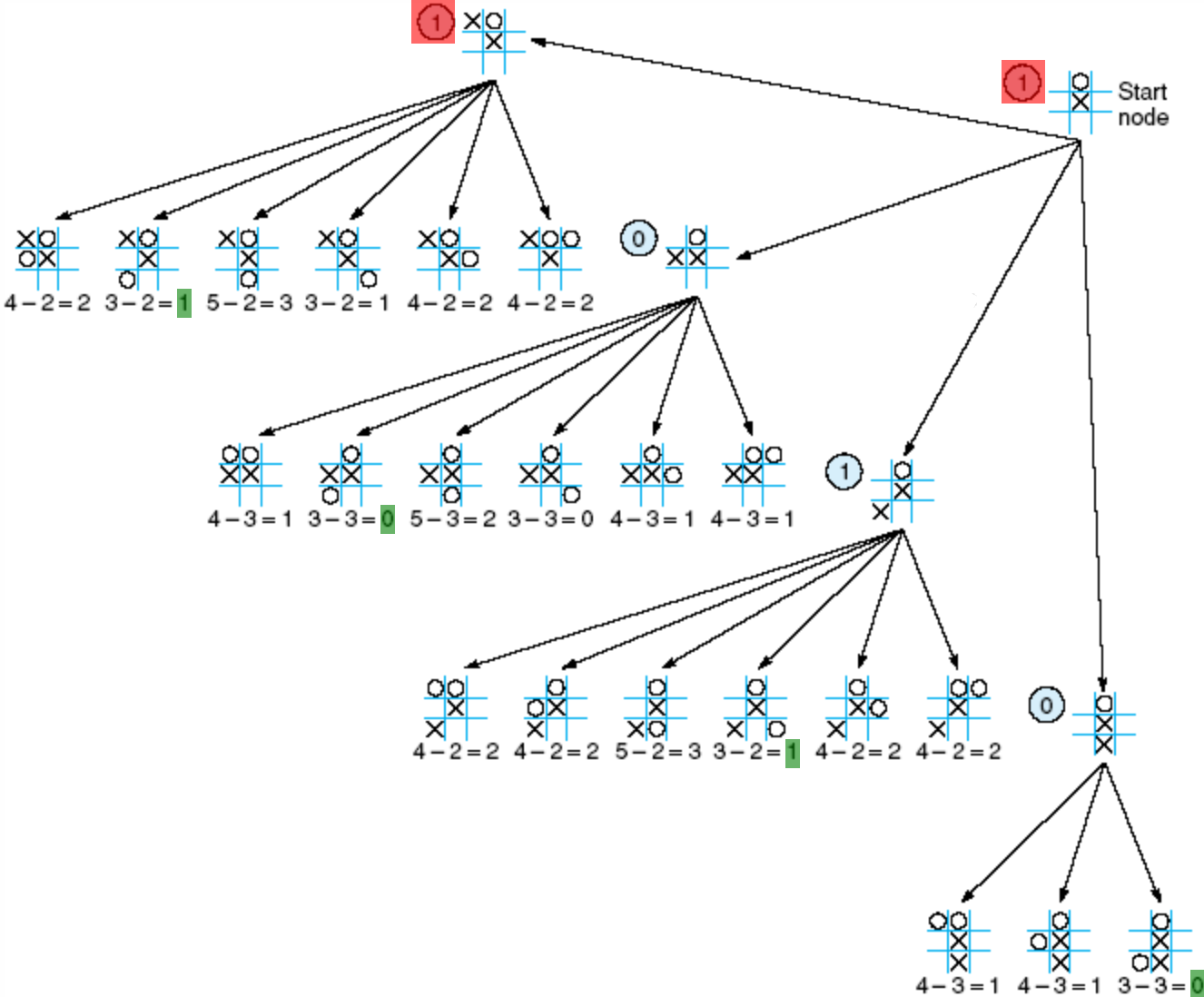
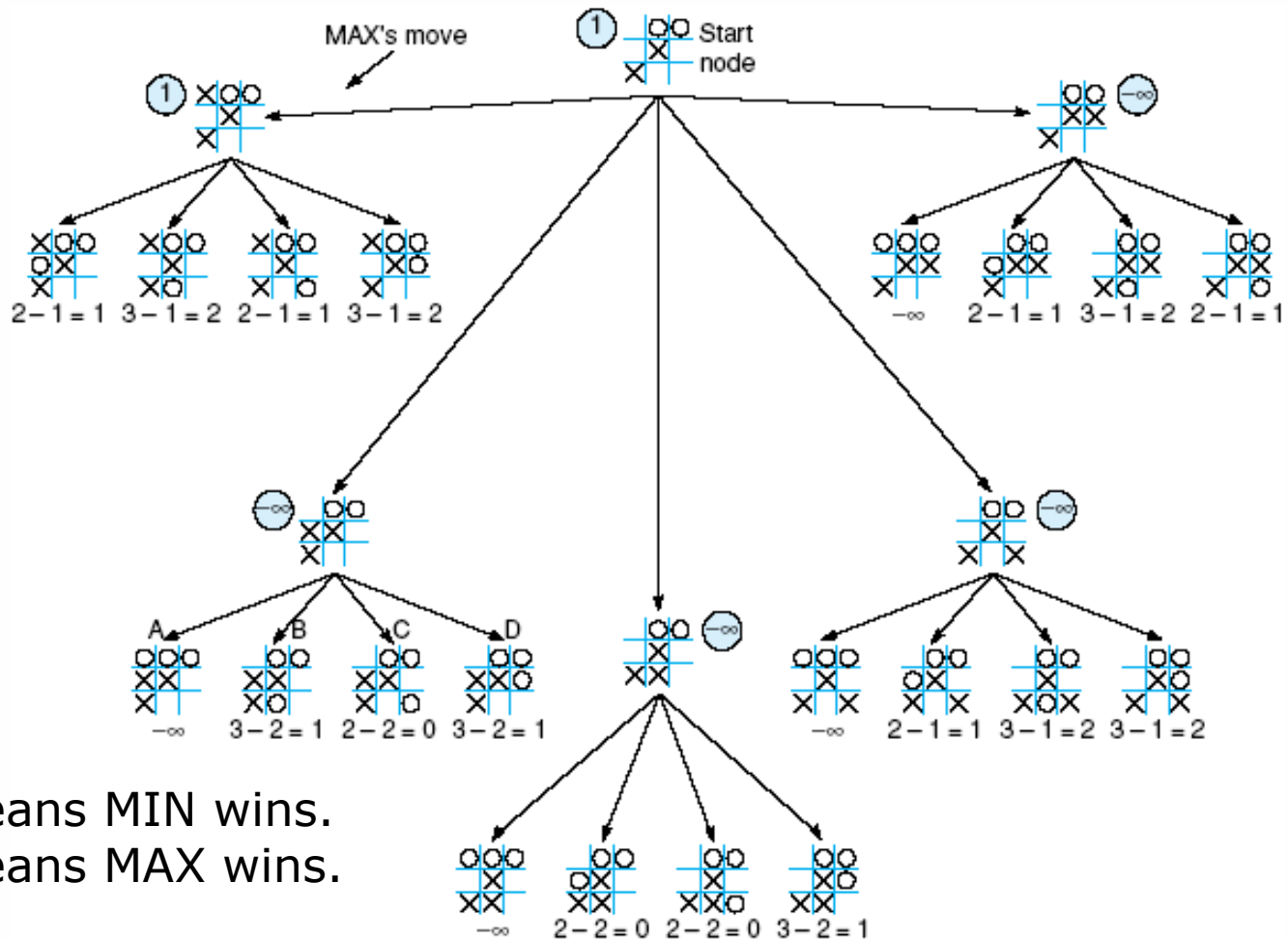
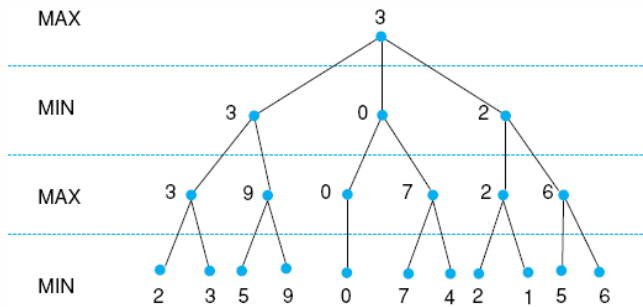


Fig. 4.24

MAX's 3rd move: look ahead analysis



Alpha-beta pruning example



Minimax without pruning

Depth-first search

Visit C, A, F,

Visit G, heuristics evaluates to 2

Visit H, heuristics evaluates to 3

Back up {2,3} to F. $\max(F)=3$

Back up to A. $\beta(A)=3$. Temporary $\min(A)$ is 3.

3 is the ceiling for node A's score.

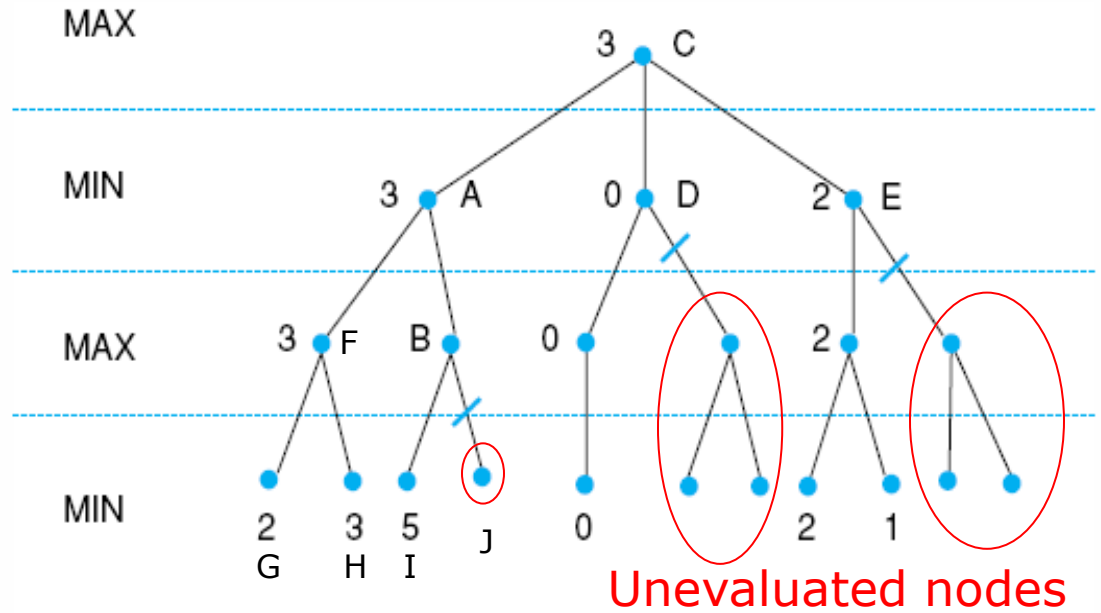
Visit B according to depth-first order.

Visit I. Evaluates to 5.

$\text{Max}(B) \geq 5$. $\alpha(B)=5$.

It does not matter what the value of J is, $\min(A)=3$. β -prune J.

Alpha-beta pruning improves search efficiency of minimax without sacrificing accuracy.



Unevaluated nodes

A has $\beta = 3$ (A will be no larger than 3)

B is β pruned, since $5 > 3$

C has $\alpha = 3$ (C will be no smaller than 3)

D is α pruned, since $0 < 3$

E is α pruned, since $2 < 3$

C is 3

Alpha-beta pruning

$\min(A) = -1$

$\max(S)$ must ≥ -1

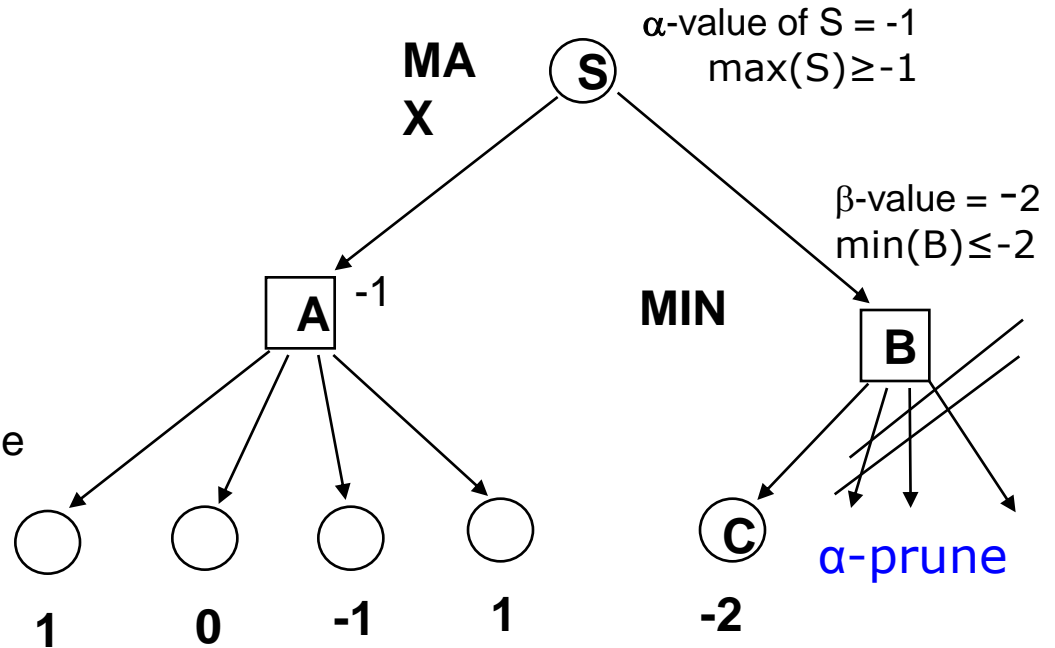
Lower bound: α -value of $S = -1$

$h(C) = -2$

$\min(B) \leq -2$

Upper bound: β -value = -2

Final value of B can never exceed current value of S . We can prune the other children of B .



- Proceed in a depth-first fashion in the n -ply look-ahead search tree.
- Find the score for the top of this tree.
- During the search, creates two values alpha and beta
- α -value associated with MAX can never decrease
 - MAX's eventually score is at least as good as the current α -value
- β -value associated with of MIN can never increase
 - MIN's eventually score is at least as good as the current β -value
- **α -prune**: β -value \leq α -value of a MAX ancestor
- β -prune: Any MAX node having α -value \leq β -value of any MIN ancestor

Conclusion

- Minimax search is designed for the adversarial search space, MAX vs MIN.
- Before MAX makes a move, he looks n plies ahead in a DFS manner.
- Apply heuristic function for the states at the end of the look ahead level.
- Propagate these values back up to the current state.
 - Use alpha-beta pruning to cut bad branches (moves) in the game tree to improve search time.
- Choose the move that maximizes the current node score.
- Then it is MIN's turn of doing similar look-ahead and pruning to decide his move.
- The players alternate until game over.